Fluid Simulation using the Lattice Boltzmann Method on the Cell Processor

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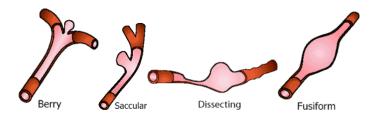
University Erlangen-Nuremberg – System Simulation

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Aneurysms

- dilatation (local balooning) of the vessel
- localized mostly at larger arteries in soft tissue (e.g. aorta, brain)
- 3-5% of all people in Germany have an aneurysm
- if an aneurysm ruptures 33% of the patients are dying and 33% suffer a physical handicap



Problem Description

Goals

blood flow simulation

- to help in understanding the development of aneurysms
- to support planning of therapy

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blood flow simulation

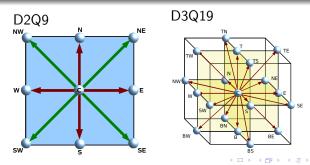
- to help in understanding the development of aneurysms
- to support planning of therapy

Challanges

- current imaging techniques (CT, MRI, Angiography) result in data sets of 512³ and more
- long run times on desktop PCs and workstations

Lattice Boltzmann Method

- fluid is described as collection of particles
- domain is divided into quadratic (2D) or cubical (3D) lattices
- every cell (lattice) has a set of particle velocity distribution functions
- simple timestep consists of two steps
 - 1 stream "move" of distribution to fluid neighbors
 - 2 collide "particle collision" leads to new distribution functions
- cellular automata, interaction only with neighbors

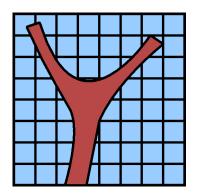


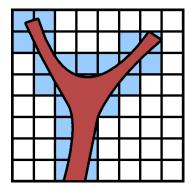
Outline

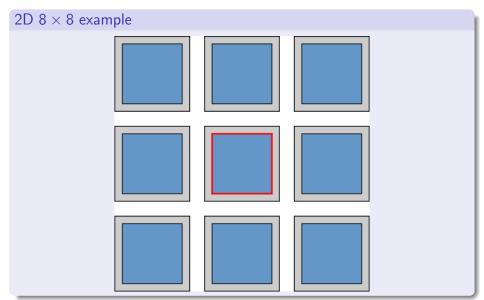
- memory layout
 - domain splitting
 - parallelization and data exchange
- SIMDization of stream step
- overview of execution on SPU
- performance
 - Rev. B Cell Blade
 - ► PS3
- outlook

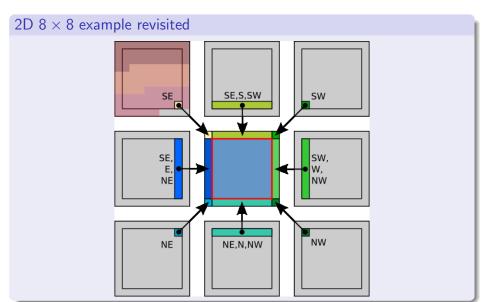
Domain Splitting

Divide whole domain into equally sized blocks ($8 \times 8 \times 8$) and only allocate and calculate blocks including fluid cells.









3D case: D3Q19

- exchange of 18 distribution functions with 18 neighbors
- 30 complete planes
 - with the six "main" neighbors (N,S,E,W,T,B)
 - five distribution functions each
- with the other 12 neighbors a single line each

3D case: D3Q19

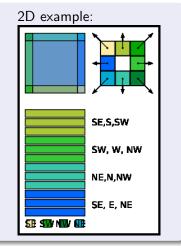
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Problems

- full set of ghost cells nearly doubles the required memory in 3D $(10^3 \text{ vs. } 8^3)$
- only few data of cache lines transferred is actually used
- scatter / gather of single values is ultra-annoying on SPUs
- lot of space wasted or badly aligned for SIMD
- parallelization will either need two versions of each block or complex memory organization to reuse data

Our Data Structure

- box' data structure contains
 - copies of outer cells to be exchanged
 - pointers to the neighbors' copies
- copies are re-ordered in SPU and stored contiguously with a single DMA transfer
- getting all necessary data from a neigboring box needs a single DMA transfer
- actually, two copy (halo) and two pointer structures (remote halo) are used
- usage of "remote halo" data is tricky



Parallelization

single node / interleaved memory

- data for all blocks necessary is allocated in one piece
- all SPU threads synchronize using atomic counting (libsync)

Parallelization

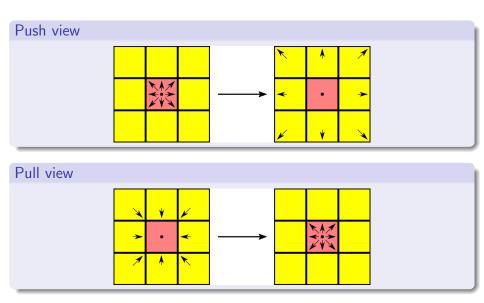
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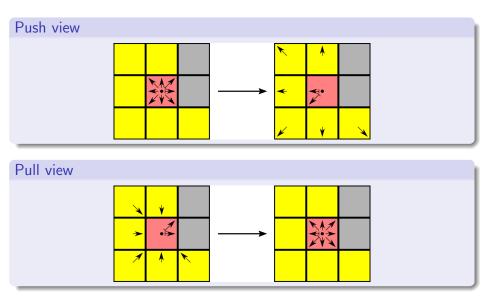
NUMA-aware

- on each node, about the same number of boxes is allocated
- SPUs on each node operate on local boxes and synchronize only locally
- remote memory access only necessary for "remote halo" transfers

Stream Step revisited



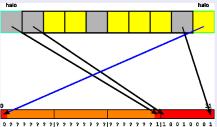
Bounce-Back - LBM No-Slip Boundary Condition



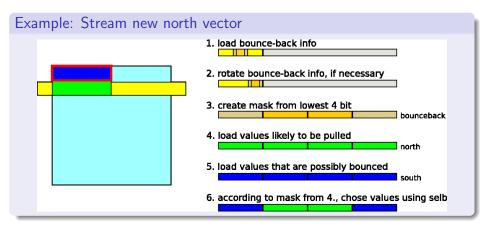
Cell Types

Cell type description

- \bullet 8 \times 8 \times 8 array of 16-bit flag fields in each block
 - source (inflow), sink (outflow)
 - experimental outflow using pressure
- type halo and type remote halo structures analogous to distribution function halos
- \bullet as geometry is currently static, all $10\times10\times10$ bounce-backs information is cached
- 100 32-bit-fields per box, describing a single line each



SIMDized Bounce-Back



SIMDized Bounce-Back

Remarks

- possibly bounced data is always aligned
- streaming must take halo data into account
 - complicated at boundaries
 - often shuffle operations to mix data from two SIMD vectors (with E or W)
- register blocking can enhance speed dramatically
 - streaming opposite directions together
 - reuse of bounce-back information
- $18 \times 8 \times 8 \times 8 = 9216$ "decisions" per block
 - scalar processing unbearably slow
 - ▶ SIMD implementations takes less than 2 clock cycles per decision
- streaming is most complicated part on SPU

Updating a Box on a SPU

Overview

- fetch...
 - distribution functions
 - ▶ LBM-cell type information
 - remote halo pointer structure
- 2 fetch...
 - halo planes from neighbors (T, B, N, S, E, W)
 - halo lines from neighbors (TS, NE, BW ...)
- set source / sink
- stream (and bounce) values (including values from halo planes)
- correct pressure outflow
- calculate collision
- prepare new halo structure from calculated values
- store box and halo data

cellIbm Performance – MFLUPS, FLUPS, FLOPS and Bandwidth

Performance numbers

MFLUPS number of lattices containing fluid updated per second

MLUPS number of lattice updates per second lattice number = box number $\times 8^3$

FLOPS the stream step usually has NO floating point operations; every collision needs 167 floating point operations (115 instructions due to 52 fused operations)

bandwidth for every lattice update at least 152B, but usually 187B and more have to be transferred

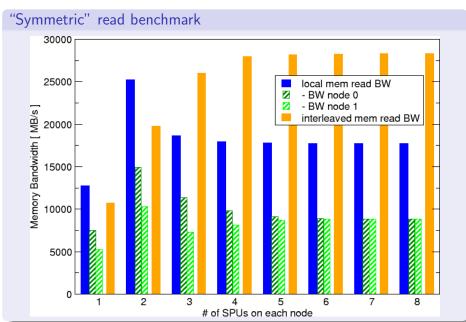
2.4 GHz Rev B Cell Blade at Böblingen

single and dual node MLUPS performance

computation	node0	parallel	parallel	
memory	node0	interleaved	distributed	
1 SPU/CPU	31	55	53	
2 SPUs/CPU	61	89	98	
3 SPUs/CPU	72	99	88	
4 SPUs/CPU	72	104	87	
96 ³ canal flow				

- strange behavior after newer kernel was installed (also using 4kB instead of 64kB pages)
- before that, maximum performance around 150 MLUPS for large domains

Problem on the Rev B Cell Blade



PS₃

"development system"

- 1 Cell Processor@3.2 GHz
- 6 SPUs available (1 disabled, 1 used from/as hypervisor)
- 256 MB Rambus XDR RAM
- 60 GB HD, gigabit ethernet
- only framebuffer graphics available
- currently running YDL 5

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Memory benchmarks on the Playstation

	bandwidth in		GB/s
page	read	write	mixed
4kB	19.7	19.4	17.6
16MB	22.6	22.6	20.3

6 SPUs doing 16k DMA transfers into a single 16MB buffer

04/11/2007

cellbm on PS3

scale up (hard)

SPUs	MFLUPS	MLUPS	
1	38	41	
2	73	78	
3	80	85	
4	80	85	
5	80	85	
96 ³ canal flow			

cellbm on PS3

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4kB pages vs. huge pages (16MB)

	MFLUPS	MLUPS	
4kB	65	72	
16MB	71	81	
48 ³ canal flow on 4 SPUs			

Outlook

More to do...

- examine and improve performance especially on Cell blade
 - influence of page size
 - kernel flaws?
 - huge pages
- MPI parallelization
 - overlapping exchange of Halo data
- examine influence of rounding mode

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Acknowledgements

- cellbm is based on Jan Götz' Master Thesis
- also lots of graphics have been shamelessly copied from him
- thanks to IBM development center at Böblingen for access to Cell Blades

The End

Thank you very much for your attention!